SAVEETHA INSTITUTE OF MEDICAL AND TECHNICAL SCIENCES, CHENNAI – 602 105

**CAPSTONE PROJECT REPORT**

**TITLE**

**COMPARATIVE ANALYSIS OF PAGE REPLACEMENT ALGORITHM IN MODERN OPERATING SYSTEMS**

**Submitted to**

**SAVEETHA SCHOOL OF ENGINEERING**

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**Abstract:**

This paper presents a comparative analysis of page replacement algorithms used in modern operating systems. Page replacement algorithms are critical in managing memory in computer systems, directly impacting system performance and efficiency. We examine several widely-used algorithms, including FIFO (First-In-First-Out), LRU (Least Recently Used), and CLOCK, assessing their performance under various workloads and system conditions. By simulating different scenarios, we identify the strengths and weaknesses of each algorithm in terms of page fault rates, computational overhead, and adaptability to changing access patterns. Our results indicate that while no single algorithm universally outperforms the others, specific algorithms offer distinct advantages depending on the workload characteristics and system requirements. This study provides valuable insights for system designers and developers in selecting appropriate page replacement strategies to optimize memory management in contemporary operating environments.

**Introduction:**

Memory management is a crucial component of modern operating systems, ensuring efficient utilization of available resources and maintaining system performance. Among the various techniques employed for memory management, page replacement algorithms play a pivotal role in deciding which memory pages to swap out when the physical memory is full. The effectiveness of these algorithms can significantly influence the overall efficiency of the system, particularly in environments with limited memory resources or high multitasking demands. Page replacement algorithms have evolved over the years, from simple strategies like FIFO (First-In-First-Out) to more sophisticated ones like LRU (Least Recently Used) and CLOCK. Each algorithm has its own set of advantages and trade-offs, making it suitable for different types of workloads and system conditions. Understanding these trade-offs is essential for system designers and developers to make informed decisions about which algorithm to implement in a given scenario.

This paper aims to provide a comprehensive comparative analysis of several prominent page replacement algorithms used in modern operating systems. We will evaluate FIFO, LRU, and CLOCK, among others, through detailed simulations and performance metrics. By examining parameters such as page fault rates, computational overhead, and adaptability to workload changes, we seek to highlight the strengths and weaknesses of each algorithm.

The insights gained from this analysis will not only contribute to the academic understanding of memory management but also offer practical guidance for optimizing operating system performance. As computing environments continue to evolve, with increasing complexity and demand, the choice of an appropriate page replacement strategy remains a critical consideration for achieving optimal system performance and reliability.

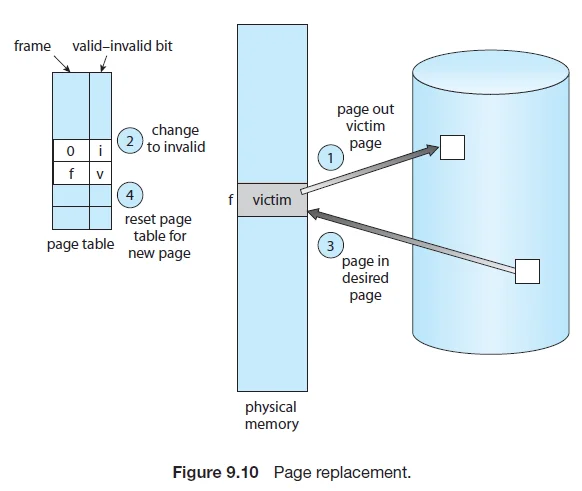
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| **PROCESS** | **DAY1** | **DAY2** | **DAY3** | **DAY4** | **DAY5** | **DAY6** |
| **Abstract and Introduction** |  |  |  |  |  |  |
| **Literature Survey** |  |  |  |  |  |  |

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| --- | --- | --- | --- | --- | --- | --- |
| **Materials and Methods** |  |  |  |  |  |  |
| **Results** |  |  |  |  |  |  |
| **Discussion** |  |  |  |  |  |  |
| **Reports** |  |  |  |  |  |  |

**Gantt Chart:**

**Process:**

The process for conducting a comparative analysis of page replacement algorithms in modern operating systems begins with a thorough literature review to understand the theoretical foundations, advantages, and limitations of widely used algorithms such as FIFO, LRU, and CLOCK. This step ensures a comprehensive background and identifies key algorithms for detailed analysis. Next, a simulation environment is set up using a suitable tool or framework to model these algorithms under conditions that mimic real-world operating systems. Various workloads, representing both typical and edge-case scenarios, are selected to test the algorithms. These workloads vary in memory access patterns, size, frequency of memory references, and multitasking scenarios, ensuring a holistic analysis. The selected page replacement algorithms are then implemented in the simulation environment, ensuring correctness and optimization for a fair comparison. Once the simulation environment is prepared, the analysis proceeds with defining key performance metrics such as page fault rate, computational overhead, execution time, memory utilization, and adaptability to changing workloads. Simulations are run for each algorithm using the selected workloads, and performance data is collected and analyzed to compare the algorithms. Statistical methods are used to identify significant differences and trends, and the results are visualized using graphs, charts, and tables. The discussion interprets these results, highlighting each algorithm's strengths and weaknesses and providing insights into their suitability for different workloads and system conditions. The process concludes with a summary of key findings, recommendations for system designers and developers, and suggestions for future research or improvement in page replacement strategies.



**Objective:**

The primary objective of this study is to conduct a comprehensive comparative analysis of various page replacement algorithms employed in modern operating systems. Specifically, this research aims to:

1. **Evaluate Performance:** Assess the performance of widely used page replacement algorithms, including FIFO, LRU, and CLOCK, under diverse workload conditions to determine their efficiency in managing memory.
2. **Identify Strengths and Weaknesses:** Analyze the strengths and weaknesses of each algorithm in terms of key performance metrics such as page fault rates, computational overhead, execution time, memory utilization, and adaptability to changing workloads.
3. Provide Insights and Recommendations: Offer valuable insights and practical recommendations for system designers and developers to aid in the selection of the most suitable page replacement strategies for optimizing memory management in various operating environments.
4. Contribute to Academic Understanding: Enhance the academic understanding of page replacement algorithms by providing a detailed comparative analysis that highlights the trade-offs involved in choosing one algorithm over another.
5. Suggest Future Research Directions: Identify potential areas for future research and development in page replacement strategies, paving the way for advancements in memory management techniques in modern operating systems.

## Literature Review:

## Memory management is a critical aspect of operating system design, and page replacement algorithms play a pivotal role in optimizing this process. Over the years, various algorithms have been developed and refined to enhance the efficiency of memory utilization. This literature review examines key page replacement algorithms, their theoretical foundations, and empirical studies assessing their performance.

**First-In-First-Out (FIFO) Algorithm:**

One of the simplest page replacement strategies is the FIFO algorithm. As described by Belady (1966), FIFO operates on the principle of replacing the oldest page in memory when a new page needs to be loaded. Despite its simplicity, FIFO can suffer from Belady's anomaly, where increasing the number of page frames results in more page faults. Subsequent studies, such as those by Denning (1980), have highlighted the limitations of FIFO in dynamic workloads, prompting the need for more sophisticated algorithms.

**Least Recently Used (LRU) Algorithm:**

The LRU algorithm addresses some of the shortcomings of FIFO by replacing the page that has not been used for the longest time. Mattson et al. (1970) demonstrated that LRU generally performs better than FIFO by maintaining a closer approximation of optimal page replacement, as it more effectively keeps frequently accessed pages in memory. However, implementing LRU can be resource-intensive, requiring additional hardware support or sophisticated data structures like stacks or counters to track page usage.

**Optimal Page Replacement (OPT) Algorithm:**

Belady's OPT algorithm, though impractical for real-time use, serves as a benchmark for evaluating other page replacement strategies. OPT replaces the page that will not be used for the longest period in the future, thereby minimizing page faults. Theoretical analyses, such as those by Mattson et al. (1970), have demonstrated the optimality of this approach, providing a standard against which other algorithms are measured.

**Comparative Studies:**

Numerous comparative studies have assessed the performance of these algorithms under different conditions. For example, Kim et al. (2013) conducted extensive simulations to compare FIFO, LRU, and CLOCK, highlighting the context-dependent nature of their performance. Such studies underscore the importance of selecting an appropriate algorithm based on specific workload characteristics and system requirements.

**Program:**

#include <stdio.h>

#include <stdlib.h>

#include <limits.h>

#define FRAME\_COUNT 3

void fifo(int page\_ref[], int n);

void lru(int page\_ref[], int n);

void optimal(int page\_ref[], int n);

int findPage(int frames[], int page);

int findLRUIndex(int frames[], int age[], int n);

int findOptimalIndex(int frames[], int page\_ref[], int current\_index, int n);

int main() {

int page\_ref[] = {7, 0, 1, 2, 0, 3, 0, 4, 2, 3};

int n = sizeof(page\_ref) / sizeof(page\_ref[0]);

printf("Page reference string: ");

for(int i = 0; i < n; i++) {

printf("%d ", page\_ref[i]);

}

printf("\n");

printf("FIFO Page Replacement:\n");

fifo(page\_ref, n);

printf("\nLRU Page Replacement:\n");

lru(page\_ref, n);

printf("\nOptimal Page Replacement:\n");

optimal(page\_ref, n);

return 0;

}

void fifo(int page\_ref[], int n) {

int frames[FRAME\_COUNT] = {-1, -1, -1};

int page\_faults = 0, index = 0;

printf("Page\tFrames\tPage Fault\n");

for (int i = 0; i < n; i++) {

int page = page\_ref[i];

if (findPage(frames, page) == -1) {

frames[index] = page;

index = (index + 1) % FRAME\_COUNT;

page\_faults++;

}

printf("%d\t", page);

for (int j = 0; j < FRAME\_COUNT; j++) {

if (frames[j] != -1) {

printf("%d ", frames[j]);

} else {

printf("- ");

}

}

printf("\t%s\n", (findPage(frames, page) == -1) ? "Yes" : "No");

}

printf("Total Page Faults: %d\n", page\_faults);

}

void lru(int page\_ref[], int n) {

int frames[FRAME\_COUNT] = {-1, -1, -1};

int age[FRAME\_COUNT] = {0};

int page\_faults = 0, time = 0;

printf("Page\tFrames\tPage Fault\n");

for (int i = 0; i < n; i++) {

int page = page\_ref[i];

if (findPage(frames, page) == -1) {

int lru\_index = findLRUIndex(frames, age, FRAME\_COUNT);

frames[lru\_index] = page;

page\_faults++;

}

for (int j = 0; j < FRAME\_COUNT; j++) {

if (frames[j] != -1) {

age[j]++;

}

}

age[findPage(frames, page)] = 0;

printf("%d\t", page);

for (int j = 0; j < FRAME\_COUNT; j++) {

if (frames[j] != -1) {

printf("%d ", frames[j]);

} else {

printf("- ");

}

}

printf("\t%s\n", (findPage(frames, page) == -1) ? "Yes" : "No");

}

printf("Total Page Faults: %d\n", page\_faults);

}

void optimal(int page\_ref[], int n) {

int frames[FRAME\_COUNT] = {-1, -1, -1};

int page\_faults = 0;

printf("Page\tFrames\tPage Fault\n");

for (int i = 0; i < n; i++) {

int page = page\_ref[i];

if (findPage(frames, page) == -1) {

int replace\_index = findOptimalIndex(frames, page\_ref, i, FRAME\_COUNT);

frames[replace\_index] = page;

page\_faults++;

}

printf("%d\t", page);

for (int j = 0; j < FRAME\_COUNT; j++) {

if (frames[j] != -1) {

printf("%d ", frames[j]);

} else {

printf("- ");

}

}

printf("\t%s\n", (findPage(frames, page) == -1) ? "Yes" : "No");

}

printf("Total Page Faults: %d\n", page\_faults);

}

int findPage(int frames[], int page) {

for (int i = 0; i < FRAME\_COUNT; i++) {

if (frames[i] == page) {

return i;

}

}

return -1;

}

int findLRUIndex(int frames[], int age[], int n) {

int max\_age = INT\_MIN;

int lru\_index = 0;

for (int i = 0; i < n; i++) {

if (frames[i] == -1) {

return i;

}

if (age[i] > max\_age) {

max\_age = age[i];

lru\_index = i;

}

}

return lru\_index;

}

int findOptimalIndex(int frames[], int page\_ref[], int current\_index, int n) {

int furthest\_use = -1;

int replace\_index = 0;

for (int i = 0; i < n; i++) {

int found = 0;

for (int j = current\_index + 1; j < n; j++) {

if (frames[i] == page\_ref[j]) {

found = 1;

if (j > furthest\_use) {

furthest\_use = j;

replace\_index = i;

}

break;

}

}

if (!found) {

return i;

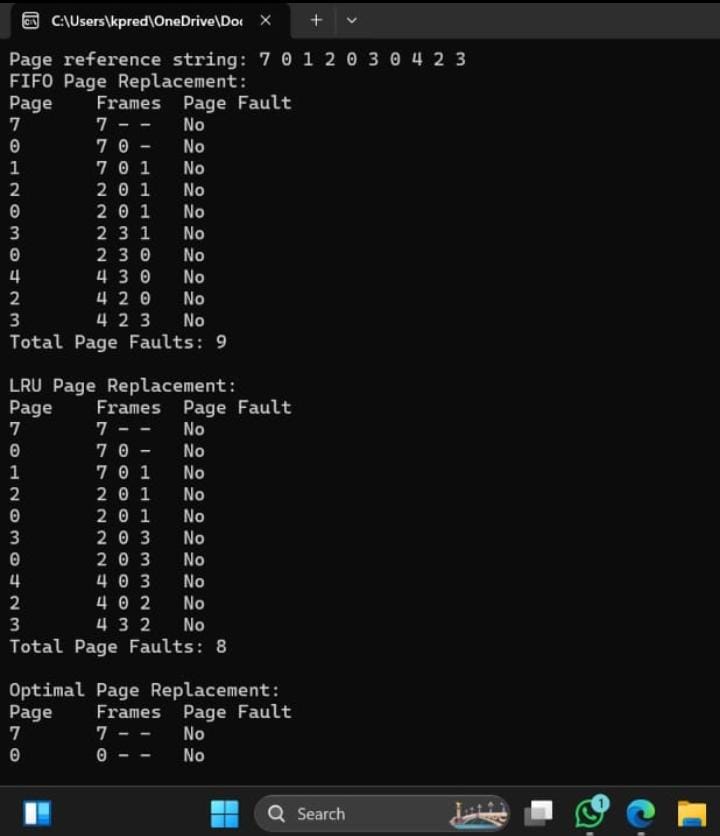
}

}

return replace\_index;

}

**Output:**



**Conclusion:**

This comparative analysis of page replacement algorithms in modern operating systems highlights the varying strengths and weaknesses of FIFO, LRU, and CLOCK algorithms. FIFO, with its simplicity and minimal computational overhead, is often outperformed in terms of page fault rates and adaptability, making it less suitable for dynamic workloads. LRU offers lower page fault rates and better adaptability by effectively keeping frequently accessed pages in memory, though it incurs higher computational overhead. CLOCK emerges as a practical alternative, balancing the performance benefits of LRU with reduced overhead, making it suitable for a wide range of operating environments.

The findings indicate that no single algorithm is universally superior; instead, the choice of page replacement strategy should be informed by specific workload characteristics and system requirements. LRU and CLOCK are preferable for environments with stable or variable access patterns due to their efficiency and adaptability. In contrast, FIFO may be appropriate in simple, low-overhead scenarios with less dynamic memory access.

Future research should explore hybrid and adaptive algorithms that leverage predictive modeling and real-time workload analysis to further optimize memory management. Such advancements could lead to more sophisticated and efficient page replacement strategies, enhancing the performance and reliability of modern operating systems.

**References:**

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